NDA: NVMe CAM attachment

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BSDCan 2016



http://people.freebsd.org/~imp/talks/bsdcan2016/slides.pdf

How I Learned To Stop Worrying and Love CAM



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http://agentpalmer.com/wp-content/uploads/2015/01/Slim-Pickens-riding-the-Bomb.jpg

Netflix

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- Internet Video
- Content Distribution Network (CDN)
- Operating at Scale
- Anticipating the Future





Netflix Open Connect

- According to Sandvine, Netflix streams ~1/3 of Internet Traffic
- Netflix has own CDN (OpenConnect)
- Streams mutliple Terabits per second



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 $http://blog.streamingmedia.com/wp-content/uploads/2014/02/2013 {\tt CDNSummit-Keynote-Netflix.pdf}$



Netflix OCA Trends

- Netflix Storage Appliance (Hard Disk Drive based)
- Netflix Flash Appliance (Solid State Drive based)
- Netflix (and industry) transitioning from SSD to NVMe

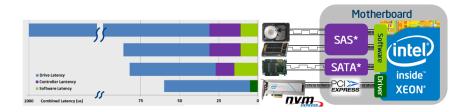


http://pcdiy.asus.com/2015/04/asus-z97-x99-motherboards-intel-750-series-nvme-ssds-all-you-need-to-know/

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Why Move To NVMe?



- 3rd Generation NVMe designs have \sim 10–15 μ s latency
- Full Bandwidth (3.9BG/s) from 4-lane PCIe Gen 3 NVMe

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• FreeBSD needs optimization (still good at \sim 30 μ s)

http://itpeernetwork.intel.com/intel-ssd-p3700-series-nvme-efficiency/



- ▶ Jim Harris of Intel wrote nvme(4) with nvd(4) disk front end
- No easy way to add I/O scheduling to nvd(4) driver
- Netflix buys cheaper drives
 - Lowers cost/GB of storage
 - More drives increases redundancy
 - Low cost drives are quirky
 - Quirkiness gets in the way of smooth, reliable performance
- CAM I/O Scheduler smooths out performance quirks





The How

- FreeBSD I/O stack overview
- CAM basics
- Structure of CAM periph (with examples from nda)
- Structure of CAM XPT (changes needed for nda)
- Structure of CAM SIM (using nvme_sim)
- Wrap up





Outline

FreeBSD I/O Stack

CAM

Code Flow Important Data Structures XPT Probe Driver Details Periph driver details XPT Details SIM drivers

Summary





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FreeBSD I/O Stack

System				
Active				
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File				
Pa	Upper ↑			
	Lower ↓			
CAM Periph Driver	mmcsd	nvd		
CAM XPT	mmcbus	nvme	NAND	
CAM SIM Driver	sdhci			
Newbus Bus Space busdma				

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After Figure 7.1 in The Design and Implementation of the FreeBSD Operating System, 2015.



FreeBSD I/O Stack

- Upper half of I/O Stack focus of VM system
 - Buffer cache
 - Memory mapped files / devices
 - Loosely coupled user actions to device action
- GEOM handles partitioning, compression, encryption
 - Filters data (compression, encryption)
 - Muxes Many to one (partitioning)
 - Muxes One to Many (striping / RAID)
 - Limited Scheduling
- CAM handles queuing and scheduling
 - Shapes flows to device
 - Limits requests to number of slots
 - Enforces rules (eg tagged vs non-tagged)
 - Multiplexes shared resources between devices

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CAM I/O Scheduler

- Written at Netflix to serve video better during "fill" periods
- Generic scheduler that allows arbitrary trade offs
- ► Gathers many real-time statistics on I/O performance
- Knows when drive has become congested

For more information please see my BSDCan 2015 I/O Scheduler talk and paper:

http://people.freebsd.org/~imp/talks/bsdcan2015/slides.pdf http://people.freebsd.org/~imp/talks/bsdcon2015/paper.pdf https://www.youtube.com/watch?v=3WqOLolj5EU

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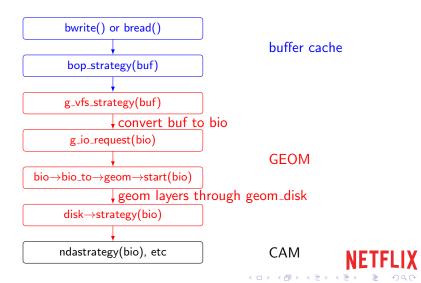
Summary





Code Flow Into CAM

File system, pager, swapper, etc





CAM Overview (Simplified)

	bio strategy() ↓	bio₋done() ↑		
Peripheral (periph)	da nda ada sa	cd ch pass ses		
Transport (XPT)	scsi ata r	nvme mmc/sd		
System Interface Module (SIM)	mpt ahci mps mp	or ahd isp nvme_sim		
	↓ hw command busdma	↑ interrupts		

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CAM Command Control Blocks (CCBs)

- Message passing mechanism of CAM
- One giant union of all possible messages
- Some commands immediate, others queued to SIM
- Completion routine to call
- Has completion status





CAM paths

- Describes nodes in the CAM device tree
- Glue that connects periph, xpt and SIM together
- All objects have one or more paths
- Allows multiple periph drivers to attach to the same device
- Includes refcounts on topology

camcontrol devlist
<Micron_M600 MU01> at scbus0 target 2 lun 0 (pass0,da0)
<Micron_M600 MU01> at scbus0 target 3 lun 0 (pass1,da1)
#

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CAM Async Notifications

- Paths register for an async notification
- Notifications queued
- Used for 'exceptional' events
 - device arrival
 - device departure
 - bus reset
- Sim gets notification to scan for devices
- XPT finds devices and gathers data
- XPT sends AC_FOUND_DEVICE and periph drivers attach

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CAM devq

- Device queuing mechanism
- One slot per slot on device
- Dynamically resizable
- Controls transactions (CCBs) sent to device
- Can be frozen for error recovery





CAM Peripheral (periph) Drivers

- Participate in device enumeration
- Take block commands via strategy function
- Convert to protocol blocks
- Send them to the SIM via the XPT
- Notifies up the stack when SIM signals completion





CAM Transport (xpt) Drivers

- Enumerates devices on transport
- Passes CCB requests from periph to SIM
- Passes CCB completions from SIM to periph
- Answers common CCBs





CAM System Interface Module (SIM) Drivers

- Not SCSI Interface Module
- Accepts protocol blocks from periph driver
- Writes CDB to host adapter
- Sets up busdma for data associated with CCB
- Signals completion of CCB when hw completion interrupt fires
- Answers CCBs about the path to the device (speed, width, mode, etc)





SIM Creation (Done In foo_attach)

- Create a devq with cam_simq_alloc
- Create a SIM with cam_sim_alloc
 - sim_action routine to receive aysnc CCBs
 - sim_poll routine for dump CCBs
 - devq
 - ▶ name / unit #
- Register each bus with xpt_bus_register
- Create a path for device enumeration with xpt_create_path

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But Where Does XPT Get Created?

- xpt_bus_register associates the xpt to the bus
- XPT_PATH_INQ CCB used to get transport type
- A giant switch statement maps the transport sub-flavors to scsi, ata, or nvme transport.
- No actual xpt object is created, just a pointer to a struct xpt_xport of function pointers.





How are periph discovered?

- Each xpt driver registers "probe" device.
- Part of the path creation process queues an AC_PATHREGISTERED notification.
- When interrupts enabled, all AC_PATHREGISTERED notifications processed.
- ► These turn into XPT_SCAN_BUS calls.
- After the probe state machine runs for each device found, the xpt layer sends AC_FOUND_DEVICE async message
- Probe devices receive these messages
- They do a XPT_PATH_INQ to discover details about the devie.

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If the details match the class of device they service, a new peripheral is added which will handle the device.



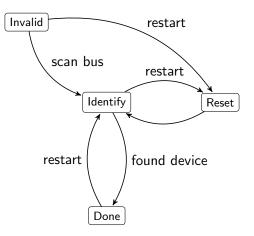
Probe state machine?

- xpt probes can't block
- xpt probes often need to send queries to the device
- State machine sends the query, when it's done the results are looked at an the next state is entered.
- For each state, a command is sent, the completion routine clocks to the next state
- Probing is done when entering the device specific done state.





NVME XPT Probe State Machine

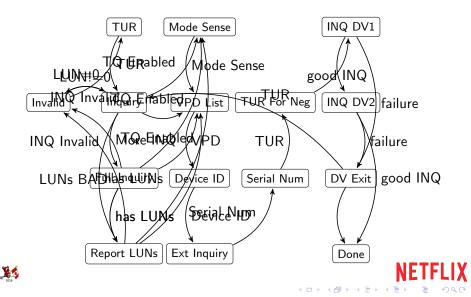


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SCSI XPT Probe State Machine



Periph driver attaching

- AC_DEVICE_FOUND sent to all devices from xpt probe
- Periph's async handler claims devices (beware: multiple can)
- Periph creates new instance of the device with cam_periph_alloc
- device's 'register' routine called
 - Allocates softc
 - Initializes I/O Scheduler
 - Matches quirks and applies them
 - Uses Inquiry or Identify Data to choose flavor of device

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- Negotiates with SIM details of the device
- Creates disk or char device
- Saves Identity information
- Registers async for interesting events
- calls xpt_schedule to get things started



Required Routines

- open Called when device is opened
- close Called on last close
- strategy Called for bio I/O
- start Called when room for work
- dump Crash dumps
- getattr Get attributes
- gone Drive has departed
- done CCB has finished





- Checks to see if there's room in devq
- If there is, it allocates a CCB and calls periph's start routine
- Can also make sure there's room in the simq for SIMs with concurrent transaction limitations beyond those of the device.





xpt_action

Pushes the I/O to XPT or SIM





xpt_done

- Finishes a CCB up and calls its completion routine
- Also calls xpt_schedule
- Requeue it if there's errors





Strategy

- System presents I/O to driver in a struct bio
- Driver queues the I/O
- Drive calls xpt_schedule to maybe do I/O





- You know you have a slot
- Must either complete CCB or submit it to SIM for I/O
- Must call xpt_schedule at the end
- Restrictions on I/O enforced here (eg, no TRIM while other I/O outstanding, etc)



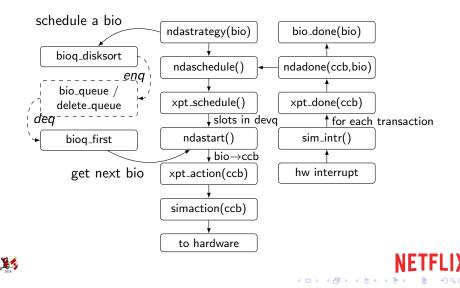


- Called by the SIM as part of xpt_done processing after it's processed the I/O
- Responsible for completing the bio up the stack
- Calls xpt_schedule since there's now a slot in drive that's opened up.





CAM I/O Code flow



SIM Routines

- simaction
- simpoll
- IRQ or Timer for completions
- created in foo_attach





simaction

- Processes the CCBs queued with xpt_action
- Queued CCBs return w/o setting the status
- Immediate CCBs do the action and set status





- Checks to see if the CCB has completed
- Called only during dumping when interrupts are disabled





sim IRQ

- Called when an I/O completes
- Finishes the CCB associated with the I/O with xpt_done





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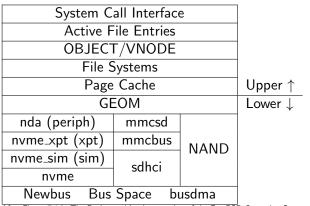
Key Points

- XPT means Transport
- SIM scans the bus for devices (explicitly, or in response to AC_PATHREGISTERED
- XPT probes device using special "probe" devices
- XPT probing state machine driven
- Once probed, XPT tells periph drivers by sending AC_FOUND_DEVICE
- periph drivers create instances based on discovered paths (may be many to 1)
- CCBs drive everything





FreeBSD I/O Stack nda World



After Figure 7.1 in The Design and Implementation of the FreeBSD Operating System, 2015.

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Questions

Questions? Comments?

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http://people.freebsd.org/~imp/talks/bsdcon2016/slides.pdf



